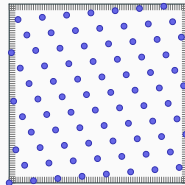
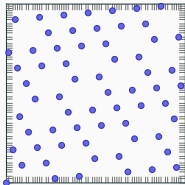
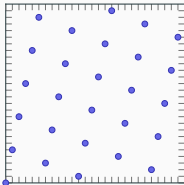


Lattice point sets and applications (part I)

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Workshop and Summer School on Applied Analysis 2023

TU Chemnitz

Chemnitz, Germany

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The plan for today

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- A light introduction to “lattice points” & “lattice rules”.
- Usage for numerical integration of “periodic” functions.
- Analysis of the error.
- Some words on function spaces and the worst-case error.
- Some Julia code to demonstrate things. . .

Lattice “points” or lattice “rules”?

Lattice rule = equal weight cubature using lattice points

For $f \in \mathcal{H}_\alpha$ approximate the d -dimensional integral

$$I(f) := \int_{[0,1]^d} f(\mathbf{x}) \, d\mathbf{x}$$

by an n -point lattice rule with generating vector $\mathbf{z} \in \mathbb{Z}_n^d$

$$Q_{n,\mathbf{z}}(f) := \frac{1}{n} \sum_{k \in \mathbb{Z}_n} f\left(\frac{\mathbf{z}k \bmod n}{n}\right).$$

Worst-case error for $f \in \mathcal{H}_\alpha$ for a given algorithm Q_n (e.g. $Q_{n,\mathbf{z}}$):

$$e^{\det}(Q_n, \mathcal{H}_\alpha) := \sup_{\substack{f \in \mathcal{H}_\alpha \\ \|f\|_\alpha \leq 1}} |I(f) - Q_n(f)|.$$

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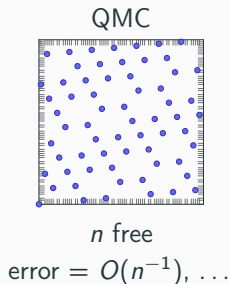
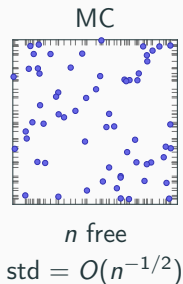
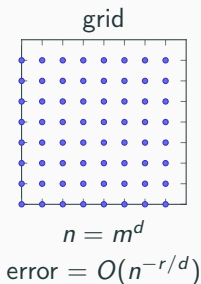
\rightsquigarrow For good lattice rule $Q_{n,\mathbf{z}}$ converges like $n^{-\alpha} \|f\|_\alpha$.

Optimal. Bakhvalov. Matching upper and lower bounds (mod logs).

“Monte Carlo type” methods: $\frac{1}{n} \sum_{k=1}^n f(\mathbf{x}_k)$

What kind of cubature/quadrature method to use for d large?

- A **product of classical quadrature rules?** (Product of weights!)
→ $n = m^d \Rightarrow$ The curse “by construction”!
- The **plain Monte Carlo method:** $\mathbf{x}_k \sim U[0, 1]^d$.
→ Free to choose n .
- **Quasi-Monte Carlo methods:** using some algebraic structure.
→ Free to choose n .



Korobov space of dominating mixed smoothness $\alpha > 0$:

$$\mathcal{H}_\alpha := \left\{ f \in L_2([0, 1]^d) : \|f\|_\alpha^2 := \sum_{\mathbf{h} \in \mathbb{Z}^d} r_\alpha^2(\mathbf{h}) |\hat{f}(\mathbf{h})|^2 < \infty \right\},$$

with

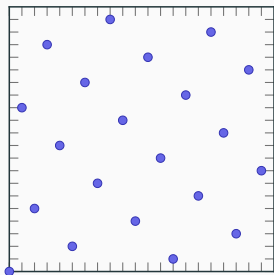
$$r_\alpha(\mathbf{h}) := \gamma_{\text{supp}(\mathbf{h})}^{-1} \prod_{j \in \text{supp}(\mathbf{h})} |h_j|^\alpha.$$

Weighted spaces: Sloan & Woźniakowski (2001),
Novak & Woźniakowski (2008, 2010, 2012), ...

More on norms tomorrow...

Example of a good lattice rule

Eg: $n = 21$ and $\mathbf{z} = (1, 13)$: Fibonacci rule: $n = F_k$, $\mathbf{z} = (1, F_{k-1})$.



Only $d = 2$, $d \geq 2$: Constructive methods for deterministic error:

Fast component-by-component (Nuyens & Cools 2006, ...)

→ Fixed vector \mathbf{z} for a given n .

(Or sequence of $n = p^m$, Cools, Kuo & Nuyens 2006).

What can we do with lattice points???

- INT: The integration problem: approximate

$$I(f) := \int_{[0,1]^d} f(\mathbf{x}) \, d\mathbf{x}.$$

- APP: The function approximation problem: find an approximation for an $f \in \mathcal{H}$ minimizing some norm.
- Collocation methods.
- Least-squares methods.
- ...

Note: If you are familiar with information based complexity (IBC): Since we use the lattice points as sample points this is the setting of [standard information](#), sometimes called Λ^{std} .

Lots of work: Korobov, Sloan, Temlyakov, Niederreiter, a lot of people in this audience...

First demo

- A (rank 1) lattice point generator (as in `Generator`).
- The “order” of the points.
- Rotated grids or grids?
- Use for numerical integration.
- Good and bad rules?

**Error for an integrand;
Worst-case error for function space**

Error for an integrand using lattice rule approximation

For $f \in \mathcal{H}_\alpha$, with $\alpha > 1/2$, or actually,
for f with abs. conv. Fourier series, “Wiener algebra”,

$$f(\mathbf{x}) = \sum_{\mathbf{h} \in \mathbb{Z}^d} \hat{f}(\mathbf{h}) e^{2\pi i \mathbf{h} \cdot \mathbf{x}}, \quad \hat{f}(\mathbf{h}) := \int_{[0,1]^d} f(\mathbf{x}) e^{-2\pi i \mathbf{h} \cdot \mathbf{x}} d\mathbf{x},$$

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we have

$$E(f) := \frac{1}{n} \sum_{k \in \mathbb{Z}_n} f\left(\frac{\mathbf{z}k \bmod n}{n}\right) - \int_{[0,1]^d} f(\mathbf{x}) d\mathbf{x} = \sum_{\substack{0 \neq \mathbf{h} \in \mathbb{Z}^d \\ \mathbf{h} \cdot \mathbf{z} \equiv 0 \pmod{n}}} \hat{f}(\mathbf{h}),$$

by the character sum for \mathbb{Z}_n , we have for $a = \mathbf{z} \cdot \mathbf{h} \in \mathbb{Z}$,

$$\frac{1}{n} \sum_{k \in \mathbb{Z}_n} \exp(2\pi i k a/n) = \mathbb{1}\{a \equiv 0 \pmod{n}\}.$$

(Show other slides with duals. . .)

Remember the definition:

Worst-case error for $f \in \mathcal{H}_\alpha$ for a given algorithm Q_n (e.g. $Q_{n,z}$):

$$e^{\det}(Q_n, \mathcal{H}_\alpha) := \sup_{\substack{f \in \mathcal{H}_\alpha \\ \|f\|_\alpha \leq 1}} |I(f) - Q_n(f)|.$$

Spaces based on series representations & Koksma–Hlawka

Assume L_2 -ONB $\{\phi_h\}_h$, $\phi_0 = 1$, $Q_n(1) = 1$, and abs. summ.

$$f(\mathbf{x}) = \sum_{\mathbf{h}} \hat{f}(\mathbf{h}) \phi_{\mathbf{h}}(\mathbf{x}), \quad \text{with} \quad \hat{f}(\mathbf{h}) := \int_{[0,1]^d} f(\mathbf{x}) \overline{\phi_{\mathbf{h}}(\mathbf{x})} d\mathbf{x},$$

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then, for $r_{\alpha,\gamma}(\mathbf{h}) > 0$ an “increasing” function,

$$\begin{aligned} |I(f) - Q_n(f)| &= \left| \sum_{\mathbf{h} \neq 0} \hat{f}(\mathbf{h}) Q_n(\phi_{\mathbf{h}}) r_{\alpha,\gamma}(\mathbf{h}) r_{\alpha,\gamma}^{-1}(\mathbf{h}) \right| \\ &\leq \left(\sum_{\mathbf{h}} \left| \hat{f}(\mathbf{h}) \right|^p r_{\alpha,\gamma}^p(\mathbf{h}) \right)^{1/p} \times \left(\sum_{\mathbf{h} \neq 0} |Q_n(\phi_{\mathbf{h}})|^q r_{\alpha,\gamma}^{-q}(\mathbf{h}) \right)^{1/q} \\ &\quad \text{norm} \qquad \qquad \qquad \times \qquad \text{worst-case error}^* . \end{aligned}$$

(See next slide.)

Worst-case error (continued...)

$$\begin{aligned} |I(f) - Q_n(f)| &= \left| \sum_{\mathbf{h} \neq 0} \hat{f}(\mathbf{h}) Q_n(\phi_{\mathbf{h}}) r_{\alpha, \gamma}(\mathbf{h}) r_{\alpha, \gamma}^{-1}(\mathbf{h}) \right| \\ &\leq \left(\sum_{\mathbf{h}} |\hat{f}(\mathbf{h})|^p r_{\alpha, \gamma}^p(\mathbf{h}) \right)^{1/p} \left(\sum_{\mathbf{h} \neq 0} |Q_n(\phi_{\mathbf{h}})|^q r_{\alpha, \gamma}^{-q}(\mathbf{h}) \right)^{1/q} \\ &\quad \text{norm} \qquad \qquad \qquad \times \qquad \text{worst-case error}^* . \end{aligned}$$

For $1 < p \leq \infty$ and compatible choices of $\phi_{\mathbf{h}}$, Q_n and $r_{\alpha, \gamma}$ we can find a “worst-case” representer $\xi(\mathbf{x})$ for which

$$|Q_n(\xi) - I(\xi)|^{1/q} = e(Q_n, \mathcal{F}_d), \quad (*)$$

independent of the particular Q_n , e.g., Fourier series and lattice rules, Walsh series and digital nets, see Nuyens (2014) and Hickernell (1998a,b).

Reproducing kernel Hilbert spaces, $p = q = 2$

Given a one-dimensional reproducing kernel $K(x, y) = \overline{K(y, x)}$.

Suppose $\mathcal{H}(K)$ is separable: $\mathcal{H}(K) = \text{span}\{\phi_h\}_h$ and $\phi_0 = 1$.

Determine the eigenvalues and eigenfunctions, and assume $\lambda_0 = 1$,

$$\int_{[0,1]} \phi(x) \overline{K(x, y)} dx = \lambda \phi(y).$$

Then

$$K(x, y) = \sum_h \frac{\phi_h(x)}{\sqrt{\lambda_h}} \frac{\overline{\phi_h(y)}}{\sqrt{\lambda_h}} = \sum_h \frac{\phi_h(x)}{\|\phi_h\|_{L_2}} \frac{\overline{\phi_h(y)}}{\|\phi_h\|_{L_2}},$$

the ϕ_h are L_2 -orthogonal, with $\|\phi_h\|_{L_2} = \sqrt{\lambda_h}$ and $\|\phi_h\|_{\mathcal{H}} = 1$, with

$$\langle f, g \rangle_{\mathcal{H}} = \sum_h \lambda_h \hat{f}(h) \overline{\hat{g}(h)}, \quad \|f\|_{\mathcal{H}}^2 = \sum_h \lambda_h |\hat{f}(h)|^2.$$

Multivariate weighted reproducing kernel Hilbert space

Use the one-dimensional space as building block for d dimensions by taking weighted tensor products (tensor product basis):

$$\begin{aligned} K(\mathbf{x}, \mathbf{y}) &= \sum_{\mathbf{u} \subseteq \{1, \dots, d\}} \gamma_{\mathbf{u}} \prod_{j \in \mathbf{u}} K(x_j, y_j) = \sum_{\mathbf{h}} \gamma_{\mathbf{u}(\mathbf{h})} \prod_{j=1}^d \frac{\phi_{h_j}(x_j)}{\sqrt{\lambda_{h_j}}} \frac{\overline{\phi_{h_j}(y_j)}}{\sqrt{\lambda_{h_j}}} \\ &= \sum_{\mathbf{h}} r_{\alpha, \gamma}^{-2}(\mathbf{h}) \phi_{\mathbf{h}}(\mathbf{x}) \overline{\phi_{\mathbf{h}}(\mathbf{y})}, \end{aligned}$$

with

(You could now interpret α as the decay of the eigenvalues.)

$$r_{\alpha, \gamma}^{-2}(\mathbf{h}) = \gamma_{\mathbf{u}(\mathbf{h})} \prod_{j \in \mathbf{u}} \lambda_{h_j}^{-1} = \gamma_{\mathbf{u}(\mathbf{h})} \prod_{j=1}^d \lambda_{h_j}^{-1},$$

and $\mathbf{u}(\mathbf{h}) = \{h_j : h_j \neq 0\} = \text{supp}(\mathbf{h})$. Now, with $\gamma_{\emptyset} = 1$, $Q_n(1) = 1$,

$$e^2(Q_n; \mathcal{H}) = -1 + \sum_{k, \ell=1}^n w_k w_{\ell} K(\mathbf{x}_k, \mathbf{y}_{\ell}).$$

For a shift-invariant space and lattice rule

For a shift-invariant space we have

$$K(\mathbf{x}, \mathbf{y}) = K(\mathbf{x} - \mathbf{y}, 0)$$

and for a lattice rule we have

$$\mathbf{x}_k - \mathbf{x}_{k'} = \mathbf{x}_{k-k' \bmod n},$$

all on the torus $[0, 1)^d$.

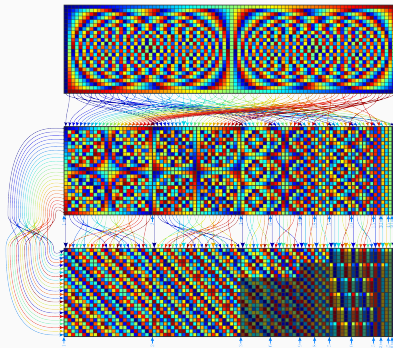
Hence:

$$\begin{aligned} e^2(Q_{n,z}; \mathcal{H}) &= -1 + \sum_{k,\ell=1}^n w_k w_\ell K(\mathbf{x}_k, \mathbf{y}_\ell) \\ &= -1 + \sum_{\ell=1}^n \frac{1}{n} \sum_{k=1}^n \frac{1}{n} K(\mathbf{x}_{k-\ell \bmod n}, 0) \\ &= -1 + \frac{1}{n} \sum_{k=1}^n K(\mathbf{x}_k, 0). \end{aligned}$$

Fast component-by-component constructions

Construction of lattice rules and polynomial lattice rules

Point sets constructed for weighted spaces using **fast component-by-component constructions** using number theoretic transforms.



See <https://www.cs.kuleuven.be/~dirkn/qmc4pde/> and <https://www.cs.kuleuven.be/~dirkn/fast-cbc/>.

See, e.g., N. & Cools (2006a,2006b), Cools, Kuo, & N. (2006), Dick, Kuo, Le Gia, N. & Schwab (2014), N. (2014), Kuo & N. (2016), ...

Variations and speedups by: Gantner, Kritzer, Laimer, Leobacher, Pillichshammer, Schwab, ... New methods: Ebert, Kritzer, N., Osisiogu (2021), Kuo, N., Wilkes (2023), N., Wilkes (2023), ...

Point generators

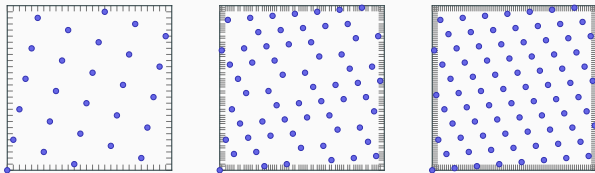
- Matlab/Octave: procedural generators like Matlab's `rand`:
 - `latticeseq_b2.m`: radical inverse lattice sequence generator,
 - `digitalseq_b2g.m`: gray coded radical inverse digital sequence generator (incl. higher-order, max 53 bit).
- Python: iterator classes, which can be used as standalone point generators from the command line (`__main__`):
 - `latticeseq_b2.py`: iterator based (`__iter__`), `set_state` for parallel computing,
 - `digitalseq_b2g.py`: ditto, arbitrary precision using `mpmath` if needed.
- C++: header file based implementation with driver program for the command line:
 - `latticeseq_b2.(h|cpp)`: complies to `ForwardIterator` concept, `set_state` for parallel computing,
 - `digitalseq_b2g.(h|cpp)`: ditto, max 64 bit.

Welcome to “The Magic Point Shop!”

Different flavours of quasi-Monte Carlo points to choose:

- Lattice rules.
- Lattice sequences.
- Polynomial lattice rules.
- Interlaced Sobol’ sequences (higher-order).
- Interlaced polynomial lattice rules (higher-order).

And code (C++ and Matlab) to use them. . .



Subsidiaries: **QMC4PDE**: construct points for parametrised PDEs.

Second demo

- The van der Corput sequence for $d = 1$.
- The Korobov trick.
- Estimating the error by use of standard error...

The end for today

The end for today from me

- Thanks for listening...

The end for today from me

- Thanks for listening. . .
- Please ask questions. . .

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- Now or later. . .

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Tomorrow more advanced things: weighted function spaces, function approximation, . . .